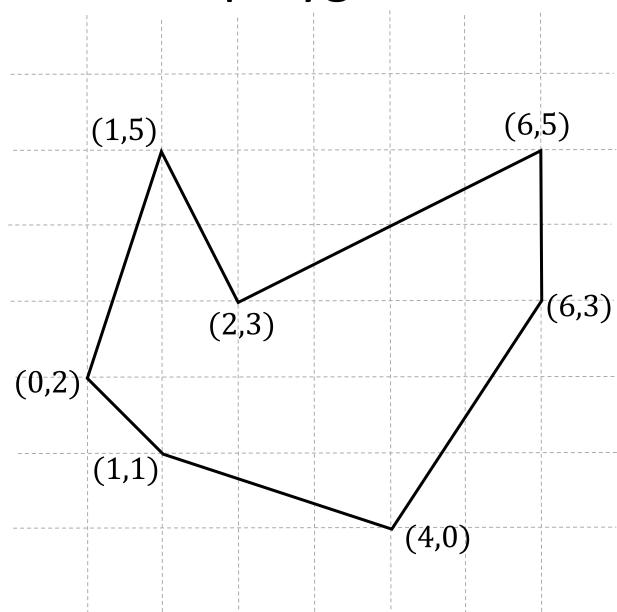
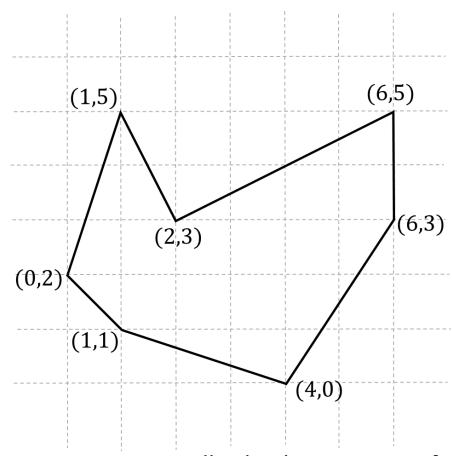
How convex is this polygon?



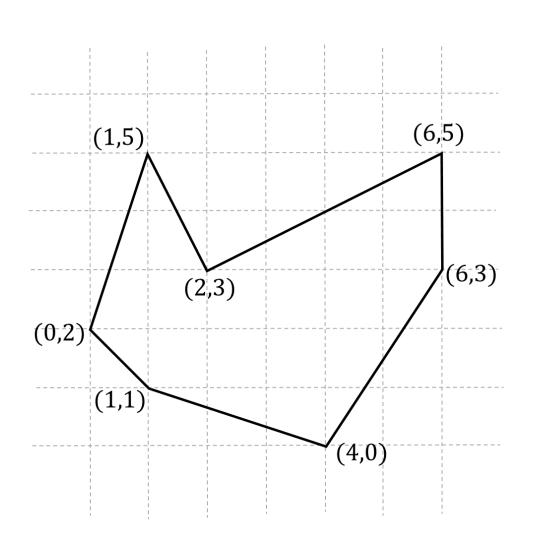
Is this polygon convex?



How can we systematically check every pair of points? A **certificate of non-convexity** is easy to describe — what about a certificate of convexity? Is this even **algorithmic**?

For any two points on the boundary, does the straight line joining them remain interior to the polygon? Maybe... Or maybe not... Is this a certificate of convexity?

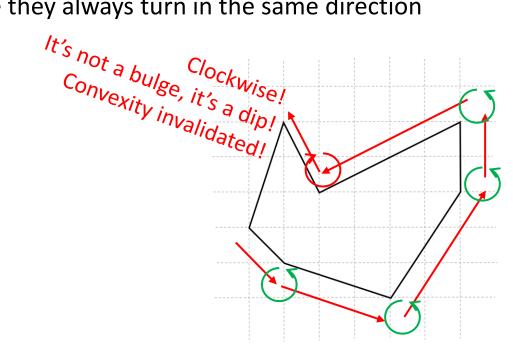
Suppose our polygon is given as a set of consecutive edge vectors



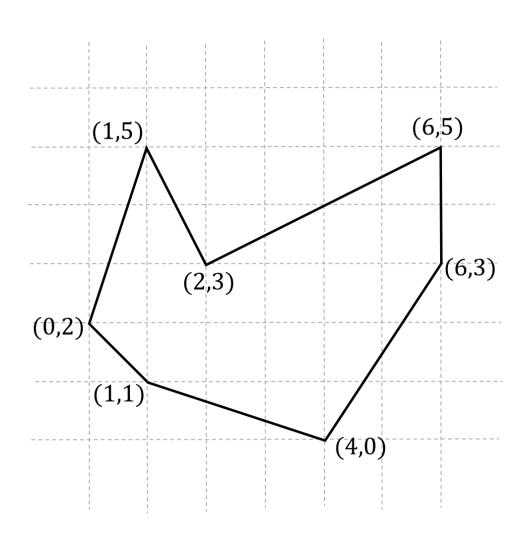
Or, same thing, a list of coordinates of vertices.

The edge from (a, b) to (c, d) is the direction vector -(a, b) + (c, d)

We will follow consecutive direction vectors and make sure they always turn in the same direction



Cross product and direction of turn



Suppose $a\mathbf{i} + b\mathbf{j} + c\mathbf{k}$ and $d\mathbf{i} + e\mathbf{j} + f\mathbf{k}$ are two vectors in three dimensions. The cross product

$$a\mathbf{i} + b\mathbf{j} + c\mathbf{k} \times d\mathbf{i} + e\mathbf{j} + f\mathbf{k}$$

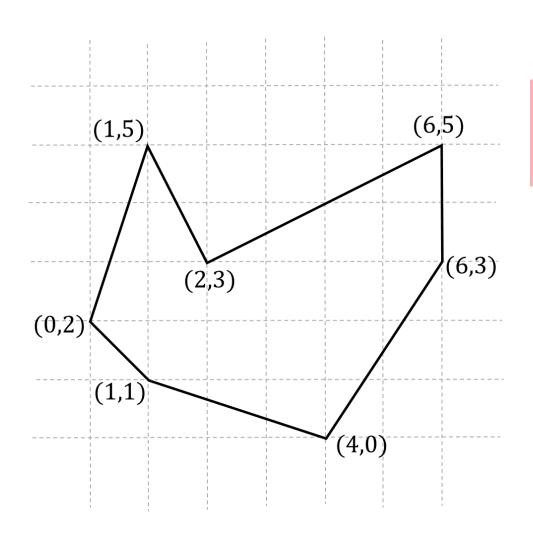
is a vector in the direction perpendicular to their common plane. It is conveniently calculated as a determinant:

It is positive if our two vectors follow each other counterclockwise and negative otherwise.

If we are in the plane then c = f = 0 and we just have $a\mathbf{i} + b\mathbf{j} \times d\mathbf{i} + e\mathbf{j} = (ae - bd)\mathbf{k}$

and we may ignore the fact that this is a vector. If the 2×2 determinant is positive we are turning counterclockwise, otherwise we are turning clockwise.

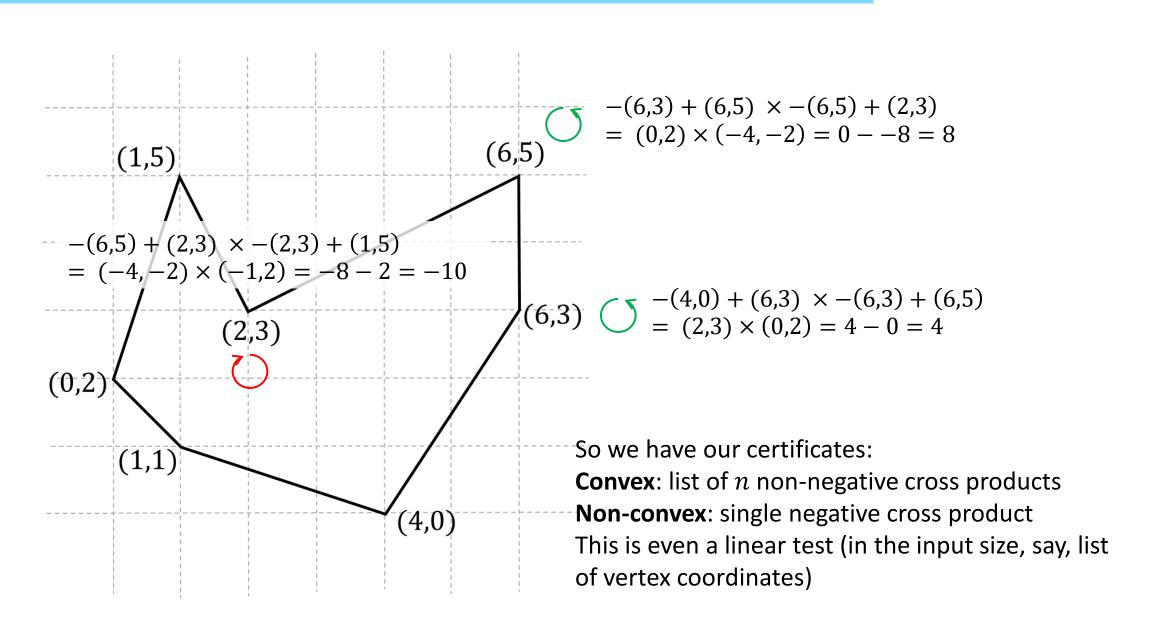
Compare to Shoelace formula



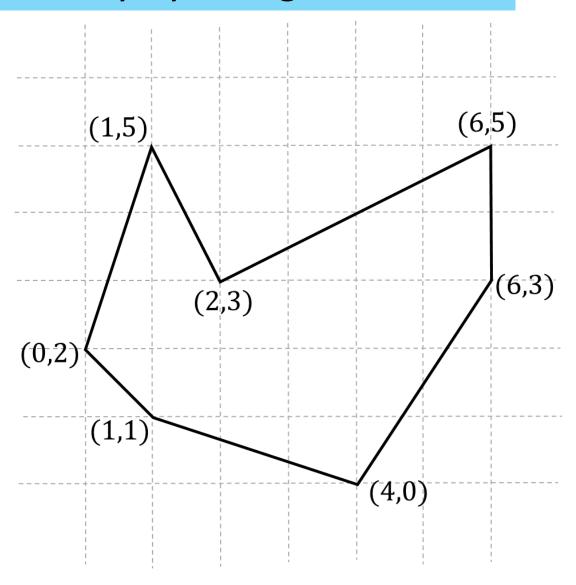
If the n vertices of a polygon are specified as position vectors $v_0, v_1, ..., v_{n-1}$, then the area of the polygon is half the sum of the cross poducts: $v_i \times v_{i+1}$, i = 0, ..., n-1.

$$2 \times \text{Area} = (0,2) \times (1,1) + (1,1) \times (4,0) + \dots + (1,5) \times (0,2).$$

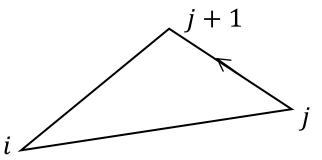
Cross product and direction of turn in the plane



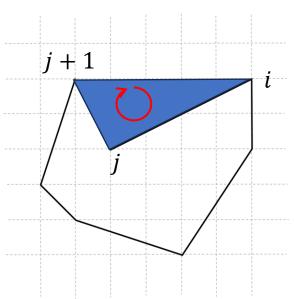
Convexity by triangulation



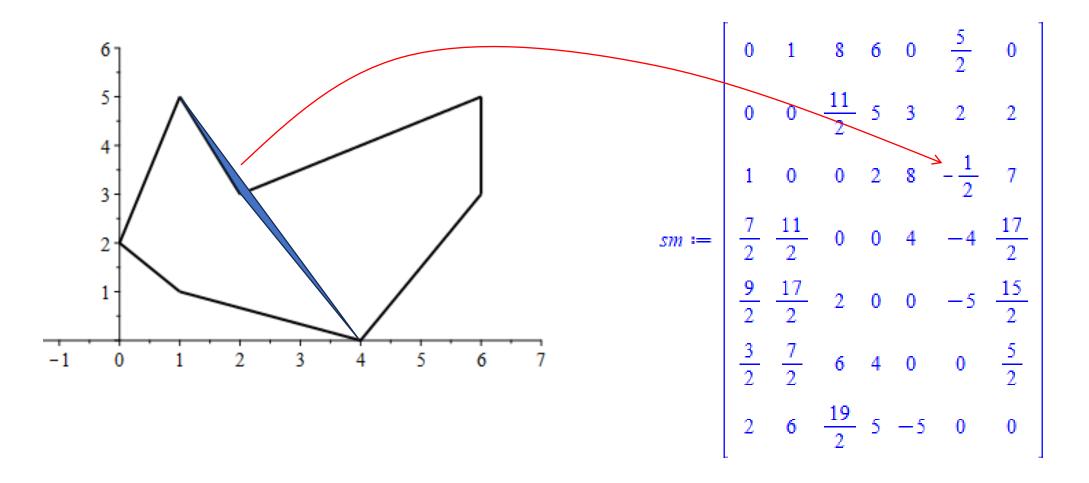
From vertex i we take the areas of all triangles subtended on opposite edges



The area is positive if the triangle has a counterclockwise orientation, otherwise it is negative



The matrix $\Delta_{i,j}$



So this is a matrix full of certificates. But it's quadratic for convexity testing isn't it? Surprisingly not.

$\Delta_{i,j}$ has rank 3

$$\begin{bmatrix} 0 & 1 & 8 & 6 & 0 & \frac{5}{2} & 0 \\ 0 & 0 & \frac{11}{2} & 5 & 3 & 2 & 2 \\ 1 & 0 & 0 & 2 & 8 & -\frac{1}{2} & 7 \\ \hline \frac{7}{2} & \frac{11}{2} & 0 & 0 & 4 & -4 & \frac{17}{2} \\ \frac{9}{2} & \frac{17}{2} & 2 & 0 & 0 & -5 & \frac{15}{2} \\ \frac{3}{2} & \frac{7}{2} & 6 & 4 & 0 & 0 & \frac{5}{2} \\ 2 & 6 & \frac{19}{2} & 5 & -5 & 0 & 0 \end{bmatrix}$$

Depends on a rather unexpected relationship between products of triangle areas:

$$\Delta_{1,2}\Delta_{0,k} - \Delta_{0,2}\Delta_{1,k} - \Delta_{0,1}\Delta_{3,k} = (\Delta_{1,2} - \Delta_{0,2} - \Delta_{0,1})\Delta_{2,k}$$

All this determined by first three rows

$$\frac{11}{2} \times \frac{5}{2} - 8 \times 2 - 1 \times -4 = \left(\frac{11}{2} - 8 - 1\right) \times -\frac{1}{2}$$

A theorem about plane triangles

Let A, B, C, D, X, Y be six points, ordered counterclockwise, in the plane.

Let ABC, ABD, ACD, BCD, be the four triangles formed on points A, B, C, D, with areas |ABC| etc.

Let Δ_A , Δ_B , Δ_C , Δ_D , be the four triangle areas formed by joining edge XY to points A, B, C, D, respectively.

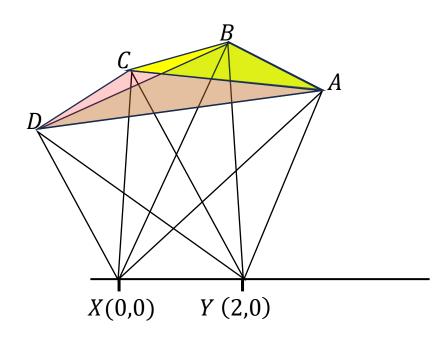
Then

$$|ABC|\Delta_D - |ABD|\Delta_C + |ACD|\Delta_B - |BCD|\Delta_A = 0.$$

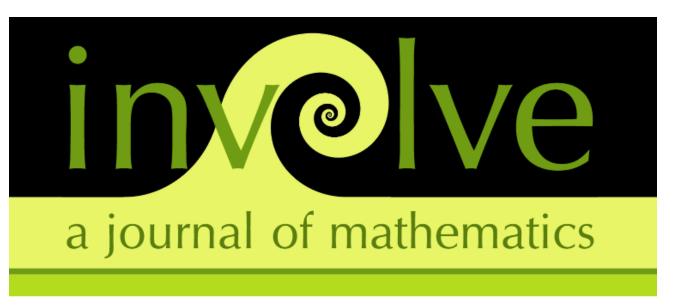
Without loss of generality, let X, Y be the points (0,0) and (2,0).

Now the areas Δ_A , etc are just the vertical coordinates of A, B, C, D, respectively.

The identity can be confirmed using the Shoelace formula.



Bisection envelopes (polygons)

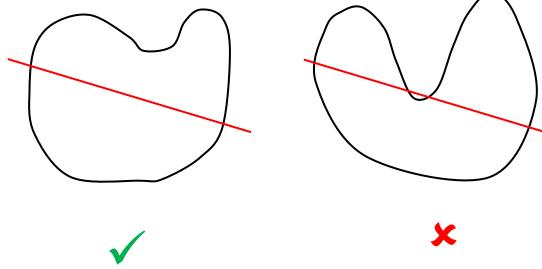


Bisection envelopes

Noah Fechtor-Pradines

Involve, Vol. 8 (2015) 307-328

Bisection-convex: any bisecting straight line intersects the curve in exactly two points



Strictly bisection-convex curves

We now restrict the class of curves S to be studied.

Definition 2.2. Define S and L as above. We say that S is *bisection convex* if for all θ , l_{θ} intersects S in exactly two points. Alternatively, for every point A on S, there exists a unique point B also on S such that the line AB bisects the interior area of S.

We also create a tighter restriction.

Definition 2.3. Define S and L as before. We say that S is *strictly bisection convex* if it is bisection convex and for all θ , l_{θ} is not tangent to S. At any point where there are two tangents to S—one from each side—the l_{θ} through that point is distinct from both tangents.

Henceforth, unless otherwise stated, it is assumed that S is strictly bisection convex.

Bisection envelopes

That IVT 2-pancakes issue again...

Define $A(\theta)$ and $B(\theta)$ to be the endpoints of the bisecting chord in direction θ , with $B(\theta) = A(\theta + \pi)$. We distinguish between $A(\theta)$ and $B(\theta)$ by demanding that for each point $Q \neq A(\theta)$, $B(\theta)$ on the bisecting chord, the vector $A(\theta) - Q$ points in positive direction θ and the vector $B(\theta) - Q$ points in positive direction $\theta + \pi$.

Proposition 2.4. Assume that S is bisection convex. Then $A(\theta)$ varies continuously with θ .

Proof. First, we note that any two bisecting chords must intersect in the interior of S, for if they did not, the interior of S would be split into three regions, one of which would have zero area, which does not make sense.

From this, we have $\lim_{\epsilon \to 0} l_{\theta+\epsilon} = l_{\theta}$, as the limit of the intersection point $l_{\theta+\epsilon} \cap l_{\theta}$ is bounded. This also implies that the limit as $\epsilon \to 0$ of the distance from $A(\theta+\epsilon)$ to the intersection point $l_{\theta+\epsilon} \cap l_{\theta}$ is bounded. Therefore, the limit as $\epsilon \to 0$ of the perpendicular distance from $A(\theta+\epsilon)$ to l_{θ} is zero.

We have that $\lim_{\epsilon \to 0} A(\theta + \epsilon)$ must be a point P on l_{θ} which intersects S, where for every other point Q on the bisecting chord with direction θ , the vector P - Q points in positive direction θ . There is only one such point, $A(\theta)$; therefore,

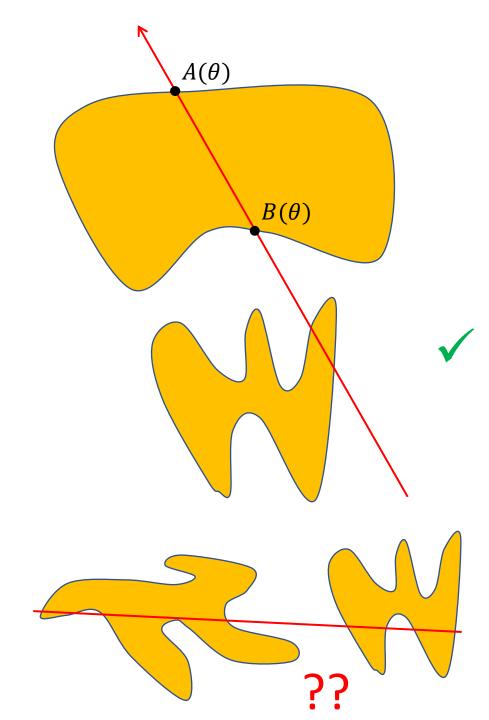
$$\lim_{\epsilon \to 0} A(\theta + \epsilon) = A(\theta),$$

and $A(\theta)$ varies continuously with θ .

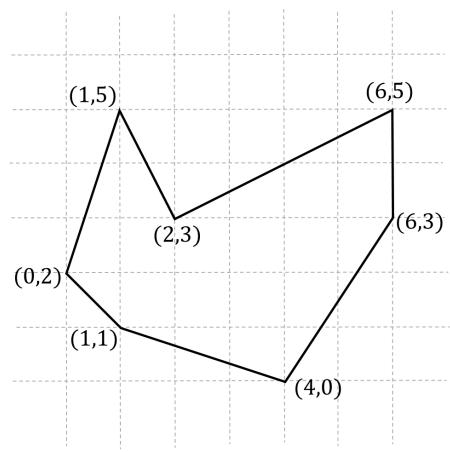
Bisection envelopes

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Is this polygon bisection-convex?



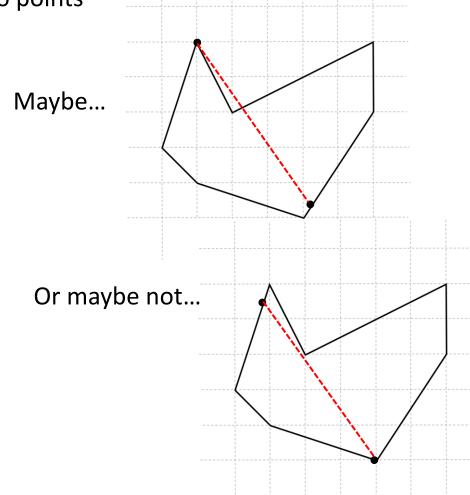
How can we systematically check every bisecting line?

A **certificate of non-bisection-convexity** is easy to describe

— what about a certificate of bisection-convexity?

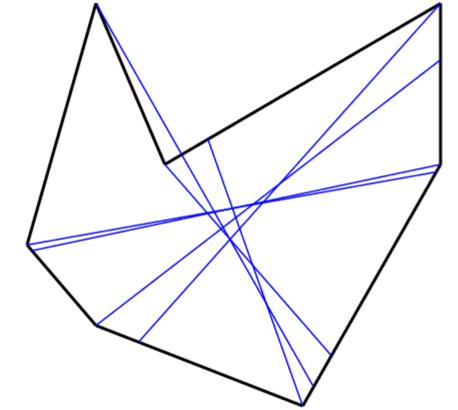
Is this even **algorithmic**?

For any straight line bisecting the polygon, does it intersect the boundary in more than two points



A characterisation

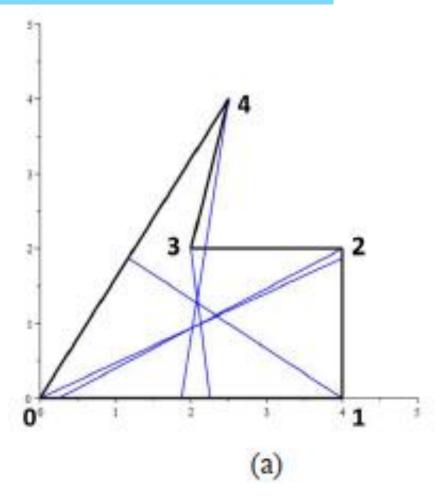
Let P be a polygon. For each vertex v of P let the unique straight line through v bisecting P be given as r_v . Then P is bisection-convex if and only if no r_v intersects the boundary of P in three or more points.



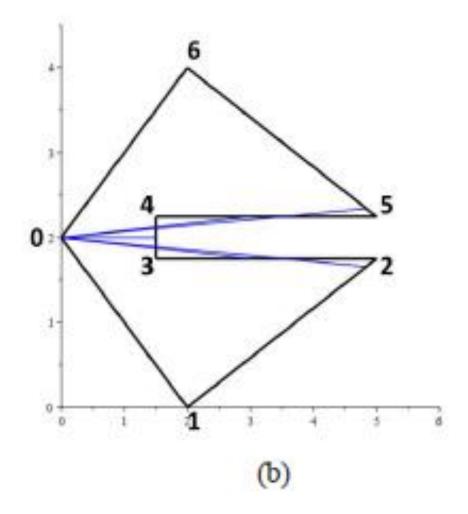
A **certificate of bisection-convexity** is a collection of n bisecting lines r_v which all lie within the boundary of P.

However, requires an effective test for a line to bisect *P*.

What is 'bisecting'?

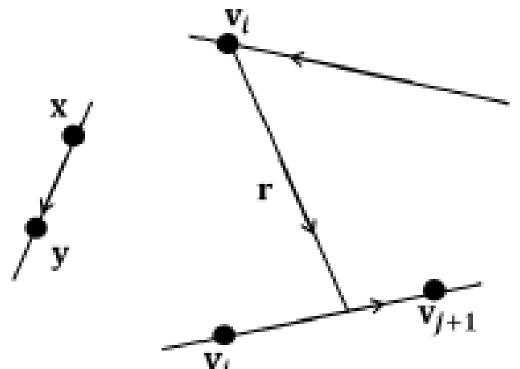


The line r_3 bisects by joining vertex 3 to an opposite edge. However, this line extends to meet the polygon boundary elsewhere.



All five lines from vertex 0 'bisect', in the sense that the two 'half' polygons joining the end-points of the lines both compute (Shoelace formula) half the area.

Is bisecting vector r crossed by edge xy?



Once again the cross product is the needed resource.

If vector r from vertex v_i crosses edge xy then $(-v_i + \mathbf{x}) \times r$ and $(-v_i + \mathbf{y}) \times r$

will have different signs.

Luckily the sequence of cross products for the polygon edges lying counterclockwise from v_i may all be calculated from the edges of the triangle area matrix $\Delta_{i,j}$

Proposition 9 Define the sequence F_i , $i \ge 0$, by

$$F_0 = 0$$
 and, for $k \ge 1$, $F_k = F_{k-1} + \Delta_{j,i+k-1} - \Delta_{i,i+k-1}$.

Then for
$$k = 1, \ldots, n-1$$
,

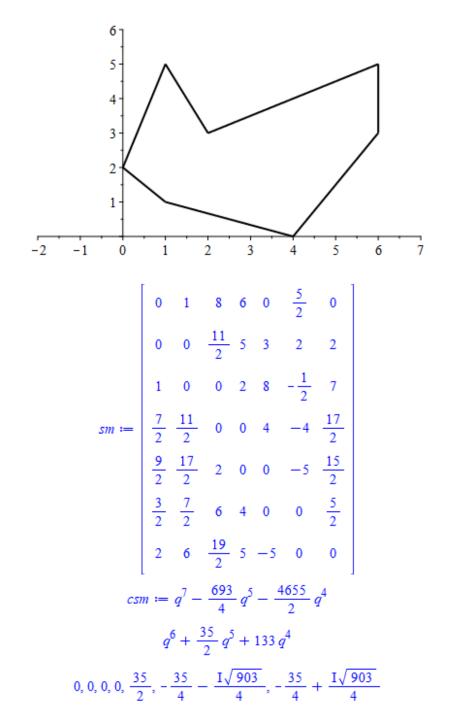
$$(-\mathbf{v}_i + \mathbf{v}_{i+k}) \times \mathbf{r} = 2r_{i,j} (\Delta_{i,j} - \Delta_{i+k,j}) + 2F_k.$$

What more can we say about $\Delta_{i,j}$?

The matrix $\Delta_{i,j}$ has rank 3, therefore n-3 zero eigenvalues. It is easy to calculate the corresponding eigenvectors using the plane triangle theorem.

The rows of $\Delta_{i,j}$ all sum to the area of the polygon, because the i-th row partitions the polygon into triangles subtended from vertex i. The area is therefore also an eigenvalue of $\Delta_{i,j}$. The corresponding eigenvector is the all-ones vector.

There remain two eigenvalues which are the (complex) roots of a quadratic polynomial. These are mysterious to me.

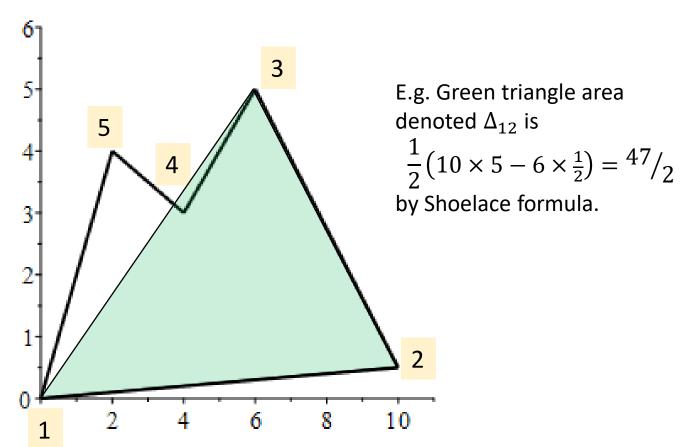


The characteristic polynomial puzzle

Characteristic polynomial of triangle areas matrix for n-vertex polygon with area P is (apparently)

$$\det(A - qI) = q^{n-3}(q - P)(q + P/2 \pm \alpha i)$$

What is α ?



Triangle areas matrix

$$(\Delta_{ij})_{i=1\dots n}$$

$$j=i+1\dots n-2+i$$

$$\begin{bmatrix} 0 & \frac{47}{2} & -1 & 5 & 0 \\ 0 & 0 & \frac{17}{2} & -\frac{1}{2} & \frac{39}{2} \\ \frac{47}{2} & 0 & 0 & -3 & 7 \\ 14 & \frac{17}{2} & 0 & 0 & 5 \\ \frac{39}{2} & 11 & -3 & 0 & 0 \end{bmatrix}$$

Char poly:
$$q\left(q-\frac{55}{2}\right)\left(q+\frac{55}{4}\pm\frac{\sqrt{5303}}{4}i\right)$$